

Some Partial Orders on Completely Regular Semigroups^X

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Abstract: The purpose of this paper is to characterize cryptogroups and normal cryptogroups by some partial order relations. It is proved that a completely regular semigroup S is a cryptogroup if and only if $S = C$ and S is a normal cryptogroup if and only if $C = S$.

Key words: completely regular semigroup; partial order; cryptogroup; normal cryptogroup

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0 Introduction and preliminaries

A completely regular semigroup S is called a cryptogroup, if Green's relation H is a congruence on S , and S/H is called a normal cryptogroup, if H is a congruence on S and S/H is a normal band.

Some partial order relations on semigroups can help us to realize the properties of semigroups and describe the structure of semigroups. For example, we know that $S = (Y; S)$ is a normal cryptogroup if and only if for any $y \in Y$, $ys = sy$ and for any $a \in S$, there exists a unique $b \in S$ such that $a = by$. And let θ be the map from S to S given by $a\theta = b$, then S is a strong semilattice of completely simple semigroups S_i and θ_i are the structure homomorphisms.

Let S be a semigroup. The binary relation C on S which was introduced by Conrad in [1] is defined as following:

$$C = \{ (a, b) \in S \times S \mid \text{for any } s \in S, asa = asb = bsa \}.$$

Burgess and Raphael proved in [2] that C is a partial order relation on S iff S is weakly separative, i. e., if $asa = asb = bsa = bsb$ for any $s \in S$, then $a = b$. In [3] Nambooripad introduced another binary relation N on a semigroup S :

$$N = \{ (a, b) \in S \times S \mid a = axa = axb = bxa, \text{ for some } x \in S \},$$

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and proved that N is a partial order on S iff S is a regular semigroup. In [4], Drazin defined a binary relation S on semigroup S :

$$S = \{ (a, b) \in S \times S \mid a^2 = ab = ba \},$$

and proved that if S is a completely regular semigroup, then S is a partial order relation on S . For a regular semigroup S , the natural partial order relation is defined as follows:

$$= \{ (a, b) \in S \times S \mid a = eb = bf, \text{ for some } e, f \in E(S) \}.$$

Drazin proved that if S is a completely regular semigroup, then $C \subseteq A \subseteq S \subseteq A \subseteq N$. It is easy to prove that $N \subseteq A$ for any regular semigroup.

A natural question is under what conditions these binary relations coincide. In this paper, we prove that if S is a completely regular semigroup, then S and N coincide if and only if S is a cryptogroup, and C and S coincide if and only if S is a normal cryptogroup.

In order to establish the main results, we need the following lemmas.

Lemma 1.^[5] Let $S = (Y; S)$ be a completely regular semigroup and \leq on Y with

() If $a \in S$, then there exists $b \in S$ such that $a \leq b$.

() If $a \in E(S)$ and $b \in S$ with $a > b$, then $b \in E(S)$.

() If $e \in E(S)$, then there exists $f \in E(S)$ such that $e \leq f$.

Unfortunately, the partial order relation S does not have the property similar to () of Lemma 1.

Example 1. Let $S = \left\{ \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}, \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}, \begin{pmatrix} 1 & 0 \\ 1 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 0 \\ -1 & 0 \end{pmatrix} \right\}$ be a matrix semigroup.

Obviously, S is a completely regular semigroup, $D = \left\{ \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}, \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} \right\}$, and $D =$

$\left\{ \begin{pmatrix} 1 & 0 \\ 1 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 0 \\ -1 & 0 \end{pmatrix} \right\}$ are two D -classes of S , and $>$. Denote $a = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} \in D$, $b =$

$\begin{pmatrix} 1 & 0 \\ 1 & 0 \end{pmatrix}$, $c = \begin{pmatrix} 1 & 0 \\ -1 & 0 \end{pmatrix} \in D$. Then $b^2 = b$, $ab = c$, $c^2 = c$, $ac = b$. So, $(c, a) \in S$ and

$(b, a) \notin S$.

Lemma 2.^[5] Let S be a completely regular semigroup. Then the following statements are equivalent:

() S is a normal cryptogroup;

() S satisfies D -majorization;

() S is a strong semilattice of completely simple semigroups;

() For all $a, x, y \in S$, $(axa)^0 = (ayxa)^0$.

Lemma 3.^[5] If S is a completely regular semigroup and μ is the largest idempotent separating congruence on S , then

() $a\mu b \iff Z a^0 = b^0$; for any $e \in E(S)$ with $e \leq a^0$, $a^{-1}ea = b^{-1}eb$,

$$() Ker\mu = \left\{ a \in S \mid \text{for any } e \in E(S) \text{ with } e^2 = a^0, ea = ae \right\}.$$

1 Main results

Proposition 1. Let S be completely regular semigroup, $a, b \in S$. Then aSb if and only if there exists $e \in E(S)$ such that $a = eb = be$.

Proof. If aSb , then $a^2 = ab = ba$ which implies $a = a^0b = ba^0$.

Conversely, if there exists $e \in E(S)$ with $a = eb = be$, then $a^2 = aeb = ab = bea = ba$, i. e., $a = a^0b = ba^0$.

Though the partial order relation S on a completely regular semigroup does not have the property similar to () of Lemma 1, we have

Proposition 2. Let $S = (Y; S)$ be a cryptogroup. Then for any $a, b \in Y$, $a, b \in S$, there exists $b \in S$ such that bSa .

Proof. By Lemma 1 and Proposition 1, for any $a \in S$, $a \in S$, there exists $e \in E(S)$ such that $e = a^0ea^0$. Since H is a congruence on S , we have $e = a^0ea^0Ha^{-1}ea$, where $a^{-1}ea \in E(S)$. So $e = a^{-1}ea$. Then $ae = aa^{-1}ea = a^0ea = ea$. Let $b = ae = ea = eae$. Then $b^0 = (eae)^0 = e$. So $b = ab^0 = b^0a$ implies $b^2 = ab = ba$, hence bSa .

Since $S \in A$ for any completely regular semigroup S , by Lemma 1 and Proposition 2, we have

Corollary 1. Let $S = (Y; S)$ be a cryptogroup, $a, b \in Y$ with $a, b \in S$, and $b \in S$. If bSa , then $b \in E(S)$.

Theorem 1. Let $S = (Y; S)$ be a completely regular semigroup. Then the partial order relations $S =$ if and only if S is a cryptogroup.

Proof. Let S be a cryptogroup. Then $\mu_S = H$.

If $a, b \in S$, $b \in a$, then $b = af = ea$ for some $e, f \in E(S)$. And $b^0 = a^0(af)^0 = (ea)^0a^0$ which implies $b^0 = a^0$. By Lemma 3, $b^0a = ab^0$. Let $b^0a = ab^0 = c$. Since $b = af = ea$, we have $b = b^0af = eab^0 = cf = ec$, which implies $b = c$. If $b \in S$ and $a \in S$ then $c = b^0a \in S$. Since the natural partial order on any completely simple semigroup is trivial, we conclude $b = c$. So, $b = ab^0 = b^0a$, and $b^2 = ab = ba$. That is, bSa . Since $S \in A$ holds for any completely regular semigroups, we get that $S =$.

If $e \in E(S)$, $e = a^0$, $b = (ea^{-1}e)^{-1}$, then $b^0 = e(ea^{-1}e)^0 = (ea^{-1}e)^0e$, i. e., $b^0 = e$ since $b^0 \in e$ and bDe . So $b = (ea^{-1}e)^{-1} = e(ea^{-1}e)^{-1} = a^0e(ea^{-1}e)^{-1} = aa^{-1}e(ea^{-1}e)^{-1}e = e(ea^{-1}e)^{-1}ea^0 = e(ea^{-1}e)^{-1}ea^{-1}a$, and $(a^{-1}e(ea^{-1}e)^{-1}e)(a^{-1}e(ea^{-1}e)^{-1}e) = a^{-1}e(ea^{-1}e)^0(ea^{-1}e)^{-1}e = a^{-1}e(ea^{-1}e)^{-1}e \in E(S)$.

Similarly, $e(ea^{-1}e)^{-1}ea^{-1} \in E(S)$. So, $b \in a$. Since $S = S$, we have $b^2 = ba = ab$, which implies $b^0a = ab^0$. Immediately, we have $ea = ae$. By Lemma 3, $Ker\mu = S$. Let aHb , i. e., $a^0 = b^0$. Then $a\mu a^0 = b^0\mu b$. So, $\mu = H$ and H is a congruence on S .

Theorem 2. Let $S = (Y; S)$ be a completely regular semigroup. Then $C = S$ if and only if S

is a normal cryptogroup.

Proof. If S is a normal cryptogroup and $a, b \in S$ with bSa , then $b^2 = ab = ba$. For any $s \in S$, from $b^2s = abs = bas$, we have $b^2sb = absb = basb$. Since S/H is a normal band, we have $bsbHb^2sb = absbHabs^2Habs$. Since $b = b^0a$, we have $bsb = (bsb)^0bsb = (bsb)^0b^0asb = (bsb)^0asb = asb$. By $b^2 = ab = ba$, we have $bsb^2 = bsab = bsba$, which induces that $bsbHbsb^2 = bsbaHb^2saHbsa$. Then $bsa = bsa(bsb)^0 = bsab^0(bsb)^0 = bsb(bsb)^0 = bsb$. Since $C \subseteq A \subseteq S$ for any completely regular semigroup^[21], we have that $C = S$.

Conversely, let $S = (Y; S)$ be a completely regular semigroup and $C = S$ on S . If $a, x, y \in S$, Then $(axya)^0 = a^0(axya)^0 = (axya)^0a^0$ which implies $(axya)^0Sa^0$. By the hypothesis, $(axya)^0Ca^0$. So, $(axya)^0(ayxa)^0(axya)^0 = a^0(ayxa)^0(axya)^0 = (axya)^0(ayxa)^0a^0$. Then $(ayxa)^0(axya)^0 = (axya)^0(ayxa)^0$. Since $(ayxa)^0$ and $(axya)^0$ lie in the same D_2 class, $(ayxa)^0 = (axya)^0$. By Lemma 2, S is a normal cryptogroup.

Proposition 3. Let S be a completely regular semigroup. Then the partial order relation $N = \dots$

Proof. If $a, b \in S$ with $a \leq b$, then there exist $e, f \in E(S)$ such that $a = be = fb$. So $a = ab^0 = ab^{-1}b$ and $a = b^0a = bb^{-1}a$. Then $ab^{-1}a = fbb^{-1}a = fb^0a = fa = fb = a$, and $a = ab^{-1}a = ab^{-1}b = bb^{-1}a$. Hence aNb . Since $N \subseteq A$ for any regular semigroup, the proposition yields.

We supply an example to explain that the condition in Proposition 3 is not necessary.

Example 2. Let $I_1 = \begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix}$, $I_2 = \begin{pmatrix} 0 & 0 \\ 1 & 0 \end{pmatrix}$, $I_3 = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}$, $I_4 = \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix}$, $I_5 = \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix}$, and $S = \{I_1, I_2, I_3, I_4, I_5\}$ be a matrix semigroup. It is easy to see that S is an inverse semigroup and is not completely regular. The binary relations $N = \{(I_1, I_2), (I_1, I_3), (I_1, I_4), (I_1, I_5)\}$ = in S .

Corollary 2. Let $S = (Y; S)$ be a completely regular semigroup. Then $S = N$ if and only if S is a cryptogroup.

Proof. The statement follows from Theorem 1 and Proposition 3.

Corollary 3. Let $S = (Y; S)$ be a completely regular semigroup. Then the partial order relations $C = S = N = \dots$ if and only if S is a normal cryptogroup.

Proof. The statement follows from Theorem 1 and Theorem 2.

Let S be a semigroup. The binary relation n should be defined as the following^[51]:

$$n = \left\{ (a, b) \in S \times S \mid a = xb = by, xa = a, \text{ for some } x, y \in S^1 \right\}$$

It is easy to verify that n is a partial order relation on S .

Proposition 4. Let S be a semigroup. Then the binary relations $N = n$ if and only if S is a regular semigroup.

Proof. If S is a regular semigroup and $a, b \in S$ with $a \leq n b$, then there exist $x, y \in S^1$ such that $a = xb = by$, $xa = a$. We can suppose that $x, y \in S$. Let $a \in V(a)$. Then $ay = xby = xa = a$, $a = aa = a = bya = a$ and $ya = a \cdot ya = ya = a$. Similarly, $aa = x \cdot aa = aa = a$.

So $a = eb = bf$ in which $e = aa^3x$, $f = ya^3a$ are idempotents of S . It is easy to verify that $a^3x, ya^3a \in V(a)$. Let $a^3 = a^3x$, $a^3 = ya^3a$. Then $a = aa^3b = ba^3a$. Let $a = a^3aa^3$. Then $aab = aa^3aa^3b = aya^3aa^3b = aa^3aa^3b = aa^3aa^3b = a$. Similarly, $a = baa = aaa$. So, aNb . That $N \setminus A \setminus n$ is trivial.

Let $N = n$ on S . Since n is a partial order relation on S , and N is a partial order relation on S if and only if S is a regular semigroup^[4], the statement holds.

Since C is always a compatible binary relation on any semigroup and C and n coincide in normal cryptogroup, we induce that the natural partial order is compatible on normal cryptogroups.

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完全正则半群上的一些偏序关系

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摘要: 用完全正则半群上的一些偏序关系刻画密码群和正规密码群. 证明了完全正则半群 S 是密码群当且仅当 $S = C$ 而 S 是正规密码群当且仅当 $C = S$.

关键词: 完全正则半群; 偏序; 密码群; 正规密码群